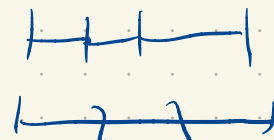


Def:  $\alpha_n: I \rightarrow S^1$   
 $\alpha_n(s) = e^{2\pi i n s}$

$\alpha = \alpha_1$                        $(\alpha \cdot \alpha) = \alpha$



Exercise:  $[\alpha]^n = [\alpha_n]$

This:  $\pi_1(S^1, 1)$  is infinite cyclic with generator  $[\alpha]$

(equivalently  $\mathbb{Z} \rightarrow \pi_1(S^1, 1)$   
 $n \mapsto [\alpha]^n$  is a group iso)

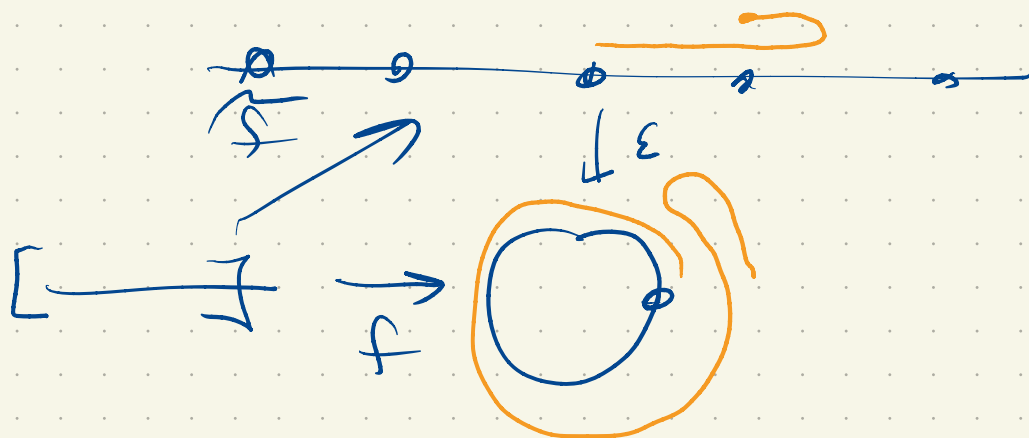
Pf: Define  $j: \mathbb{Z} \rightarrow \pi_1(S^1, 1)$  by  $j(n) = [\alpha]^n$ .

Observe  $j(n+m) = [\alpha]^{(n+m)} = [\alpha]^n [\alpha]^m = j(n) \circ j(m)$ ,

So  $j$  is a group homomorphism.

To see that  $j$  is surjective consider  $[f] \in \pi_1(S, 1)$ ,

Let  $\tilde{f}$  be a lift of  $f$  starting at 0.



Let  $n = \tilde{f}(1)$ .

Let  $H(s, t) = (1-t)\tilde{f}(s) + tns$



Observe that  $H$  is a path homotopy.

Moreover  $\epsilon \circ H$  is a path homotopy from  $f$  to  $\left( \begin{array}{c} s \mapsto \epsilon(n \circ s) \\ \epsilon \pi_1(S) \end{array} \right)$

$\alpha_n$ . So  $[f] = j(n)$ .

To establish injectivity suppose  $j(n) = 1 = [c_1]$   
 $= [\alpha_0]$ .

We need to show  $n = 0$ . Let  $H$  be  
a path homotopy from  $j(n)$  to  $[x_0]$ , i.e. from  $[\alpha_n]$  to  $[\alpha_0]$ .



By homotopy lifting we obtain

$$\tilde{H}: I \times I \rightarrow \mathbb{R} \quad \text{with} \quad \varepsilon \circ \tilde{H} = H$$

and such that  $\tilde{H}(0,0) = 0$ .

Using the fact that constants lift to constants

we obtain  $\tilde{H} = 0$  on three sides,

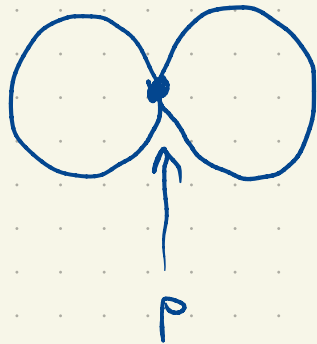
In particular  $\tilde{H}(1,0) = 0$ .



But  $\tilde{H}(s, 0)$  is a lift of  $\alpha_n$  starting  
at  $0$ . Hence  $\tilde{H}(s, 0) = ns$ . Since  $\tilde{H}(1, 0) = 0$   
we conclude  $n = 0$ .

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Fundamental Groups from precos



Let  $G_1$  and  $G_2$  be groups  $G_1 \cap G_2 = \emptyset$  for simplicity

$\prod_{\alpha \in I, 2} G_\alpha$  A word in  $G_1 \cup G_2$  is a finite tuple, possibly empty,  $(g_1, \dots, g_n)$  with  $g_i \in G_1 \cup G_2$ .

We have a product on words  $(g_1, \dots, g_n) \circ (h_1, \dots, h_m)$   
 $= (g_1, \dots, g_n, h_1, \dots, h_m)$ .

$$g, g' \in G_1$$

$$g_1, g_2, \dots, g_n$$

$$(g) (g') = (gg')$$

$$\begin{array}{ccc} & \downarrow & \\ (g, g') & \longrightarrow & (1_\alpha) \end{array}$$

$$\longrightarrow (1_\alpha)$$

## Elementary reductions

$$(g_1, \dots, g_{i-1}, 1_\alpha, g_{i+1}, \dots, g_n) \longrightarrow (g_1, \dots, g_{i-1}, g_{i+1}, \dots, g_n)$$

$$\text{If } g_i, g_{i+1} \in G_\alpha$$

$$(g_1, \dots, g_{i-1}, g_i, g_{i+1}, \dots, g_n) \longrightarrow (g_1, \dots, g_{i-1}, g_i g_{i+1}, \dots, g_n)$$

We say words  $W, W'$  are related if there is a finite sequence

$$W = W_1, W_2, \dots, W_m = W'$$

such that for each  $j$  there is an elementary reduction taking  $W_j$  to  $W_{j+1}$  or vice-versa,

Exercise: This is an equivalence relation,

Def:  $G_1 * G_2$  (the free product of  $G_1$  with  $G_2$ )

is the set of equivalence classes of words

in  $G_1 \cup G_2$  under this equiv. relation,

We define a product on  $G_1 * G_2$  by

$$[\underbrace{w_1}_{w_1}] \cdot [\underbrace{w_2}_{w_2}] = [w_1 w_2].$$

Exercise: This is well defined,

$$[(1_{G_1})] \stackrel{?}{=} [()]$$

The identity element is  $[()]$

$$\text{If } W = (g_1, \dots, g_n)$$

$$[W]^{-1} = [(g_1^{-1}, \dots, g_n^{-1})]$$

$$[W] [(g_1^{-1}, \dots, g_n^{-1})]$$

$$= [(g_1, \dots, g_n, \underbrace{g_1^{-1}, \dots, g_n^{-1}})]$$

$$= [(g_1, \dots, g_n g_n^{-1}, \dots, g_1^{-1})]$$

$$= [(g_1, \dots, \underbrace{1}_{q}, \dots, g_1^{-1})]$$

$$= [(g_1, \dots, g_{n-1}, g_{n-1}^{-1}, \dots, g_1^{-1})]$$



$$\approx [ (g, g^{-1}) ]$$

$$\approx [ (I_\alpha) ]$$

$$\approx [ () ]$$

Associativity:

$$\begin{aligned} ([w_1] [w_2]) [w_3] &= ([w_1, w_2]) [w_3] \\ &= [(w_1, w_2) w_3] \\ &= [w_1 w_2 w_3] \end{aligned}$$

If  $g_1 \in G_1$  and  $g_2 \in G_2$   $g_1 \neq 1_d$   $g_2 \neq id$ ,

$$\text{is } g_1 g_2 = g_2 g_1$$

No!

We say a word is reduced if it contains no identity elements and no two adjacent entries come from the same group.

Claim: every word is related to a unique reduced word.

Plan: I'm going to build

$$r: W \rightarrow R$$

↑  
words

↑  
reduced words

such that 1)  $r(w) = w$  if  $w$  is reduced

2)  $r(w) = r(w')$  if  $w \sim w'$

If I do this suppose  $w$  is related to  
the reduced words  $V, V'$

Then  $V' = r(V') = r(w) = r(V) = V$